**Example 101.** Let A be a  $3 \times 5$  matrix. For each of col(A), row(A) and null(A) determine d so that the space is a subspace of  $\mathbb{R}^d$ .

**Solution.** col(A) is a subspace of  $\mathbb{R}^3$ . row(A) is a subspace of  $\mathbb{R}^5$ . null(A) is a subspace of  $\mathbb{R}^5$ .

Example 102. Let  $A = \begin{bmatrix} 1 & 2 & 0 & 4 \\ 2 & 4 & 0 & 8 \\ 3 & 6 & 2 & 22 \\ 4 & 8 & 0 & 16 \end{bmatrix}$ . Find a basis for col(A), row(A) and null(A).

**Solution.** We compute an echelon form of *A*:

$\begin{bmatrix} 1 & 2 & 0 & 4 \end{bmatrix}$	$\begin{array}{c} R_2 - 2R_1 \Rightarrow R_2 \\ R_3 - 3R_1 \Rightarrow R_3 \end{array}$	1	<b>2</b>	0	4		1	<b>2</b>	0	4 ]	
2 4 0 8	$R_4 - 4R_1 \Rightarrow R_4$				0	$R_2 \Leftrightarrow R_3 $	0	0	<b>2</b>	10	
3 6 2 22	$\sim \rightarrow$	0	0	<b>2</b>	10	$\sim \rightarrow$	0	0	0	0	
4 8 0 16		0	0	0	0		0	0	0	0	

We can now use Theorem 99 to read off the bases we are interested in:

• The pivot columns are the first and third. Hence, a basis for col(A) is  $\begin{bmatrix} 1 & 0 & 0 \\ 2 & 3 \\ 4 & 0 \end{bmatrix}$ .

[Make sure you see why it would be horribly wrong to take columns from the echelon form.]

•  $\begin{bmatrix} 1 \\ 2 \\ 0 \\ 4 \end{bmatrix}$ ,  $\begin{bmatrix} 0 \\ 0 \\ 2 \\ 10 \end{bmatrix}$  form a basis for row(A).

[Note that it would be horribly wrong to take the first two rows from A.]

• The general solution to 
$$A\boldsymbol{x} = \boldsymbol{0}$$
 is  $\boldsymbol{x} = \begin{bmatrix} -2s_1 - 4s_2 \\ s_1 \\ -5s_2 \\ s_2 \end{bmatrix} = s_1 \begin{bmatrix} -2 \\ 1 \\ 0 \\ 0 \end{bmatrix} + s_2 \begin{bmatrix} -4 \\ 0 \\ -5 \\ 1 \end{bmatrix}$ .  
Hence,  $\begin{bmatrix} -2 \\ 1 \\ 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} -4 \\ 0 \\ -5 \\ 1 \end{bmatrix}$  are a basis for null(A).

Let A be  $m \times n$ , and let r be the **rank** of A, that is, r is the number of pivots.

- $\dim \operatorname{col}(A) = \dim \operatorname{row}(A) = r$
- $\dim \operatorname{null}(A) = n r$

**Example 103.** The  $4 \times 4$  matrix  $A = \begin{bmatrix} 1 & 2 & 0 & 4 \\ 2 & 4 & 0 & 8 \\ 3 & 6 & 2 & 22 \\ 4 & 8 & 0 & 16 \end{bmatrix}$  from the previous example has rank 2.

Indeed,  $\dim \operatorname{col}(A) = 2$ ,  $\dim \operatorname{row}(A) = 2$ ,  $\dim \operatorname{null}(A) = 4 - 2 = 2$  matches our computations.

**Example 104.** Let A be a  $3 \times 5$  matrix of rank 2. Determine the dimensions of col(A), row(A) and null(A).

**Solution.** dim col(A) = 2, dim row(A) = 2, dim null(A) = 5 - 2 = 3.

**Example 105.** Find a basis for col(A), row(A), null(A) with  $A = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 0 \\ 2 & 4 & 0 \end{bmatrix}$ .

Solution. For this simple matrix, we can just "see" the following (make sure you do, too!):

• A basis for col(A) is:  $\begin{bmatrix} 1\\0\\2 \end{bmatrix}$ ,  $\begin{bmatrix} 1\\0\\4 \end{bmatrix}$  Why? Because these two vectors span and are clearly independent.

Note. We would select the same basis, if we computed an echelon form of A and applied Theorem 99(a).

• A basis for row(A) is:  $\begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 2 \\ 4 \\ 0 \end{bmatrix}$  Why? Again, because these two vectors span and are clearly independent.

Note. An echelon form of A is  $\begin{bmatrix} 1 & 1 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 0 \end{bmatrix}$ . By Theorem 99(b), an alternative basis for row(A) is  $\begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 0 \\ 2 \\ 0 \\ 0 \end{bmatrix}$ . In fact, further computing the RREF, we would select as basis  $\begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}$ ,  $\begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$  (probably the nicest basis for most purposes).

• A basis for null(A) is:  $\begin{bmatrix} 0\\0\\1 \end{bmatrix}$ 

Why? We know that rank(A) = 2. Hence, dim null(A) = 3 - 2 = 1. Therefore, any nonzero vector in null(A) will be a basis for null(A). Clearly,  $\begin{bmatrix} 0 & 0 & 1 \end{bmatrix}^T$  is one such vector solving  $A\mathbf{x} = \mathbf{0}$  (why?).

## How little we actually know!

## **Q**: How fast can we solve N linear equations in N unknowns?

Estimated cost of Gaussian elimination:

 $\begin{bmatrix} \bullet & * & * & \cdots & * \\ 0 & * & * & \cdots & * \\ \vdots & \vdots & & \vdots \\ 0 & * & * & \cdots & * \end{bmatrix}$ • to create the zeros below the first pivot:  $\Rightarrow$  on the order of  $N^2$  operations • if there is N pivots total:  $\Rightarrow$  on the order of  $N \cdot N^2 = N^3$  operations

- A more careful count places the cost at  $\sim \frac{1}{3}N^3$  operations.
- For large N, it is only the  $N^3$  that matters. It says that if  $N \rightarrow 10N$  then we have to work 1000 times as hard.

That's not optimal! We can do better than Gaussian elimination:

- Strassen algorithm (1969):  $N^{\log_2 7} = N^{2.807}$
- Coppersmith–Winograd algorithm (1990):  $N^{2.375}$
- ... Stothers–Williams–Le Gall (2014):  $N^{2.373}$  (If  $N \rightarrow 10N$  then we have to work 229 times as hard.)

Is  $N^{2+(a \text{ tiny bit})}$  possible? We don't know! (People increasingly suspect so.) (Better than  $N^2$  is impossible; why?)

## Good news for applications:

• Matrices typically have lots of structure and zeros which makes solving so much faster.